

During the lecture, one key conclusion we discovered was that $\log(N!) \in \Theta(N \log N)$ and also that $N \log N \in \Theta(\log(N!))$. After proving the Big Omega bound in one direction, we then showed the Big Omega bound in the other direction, $N \log N \in \Omega(\log(N!))$. Which of the following steps were included in the derivation?

* 4 points

- ☐ $N \log N = \log(N) + \log(N) + \dots + \log(N)$ where we add $\log(N)$ N times
- ☐ $\log(N!) = \log(N) + \log(N-1) + \log(N-2) + \dots + \log(1)$
- ☐ Asymptotic analysis drops multiplicative factors and lower order terms so $N! \in \Theta(N)$
- ☐ Comparing terms after rewriting $N \log N$ and $\log(N!)$, we see that $\log(N!)$ is always less than or equal to $N \log N$

To tighten the lower bound and show that asymptotically fastest comparison sort is $\Omega(N \log N)$, we used the puppy, cat, dog problem. Which of the following statements is true about sorting and the puppy, cat, dog problem.

* 4 points

- ☐ The puppy, cat, dog problem reduces to sorting because we can use sorting to solve the puppy, cat, dog problem
- ☐ Sorting reduces to the puppy, cat, dog problem because we can use the puppy, cat, dog problem to solve sorting
- ☐ The puppy, cat, dog problem on N items requires at least $N \log N$ comparisons
- ☐ If comparison sorts could be faster than $\Omega(N \log N)$, then puppy, cat, dog could also be solved in fewer than $\Omega(N \log N)$ comparisons

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