

Select the appropriate asymptotic expression that bounds each sum defined below:

* 6 points

Not all expressions will be used.

1. $1 + 2 + 3 + \dots + N$

2. $2^1 + 2^2 + 2^3 + \dots + 2^N$

3. $3^1 + 3^2 + 3^3 + \dots + 3^N$

	$\theta(\log N)$	$\theta(N)$	$\theta(N^2)$	$\theta(2^N)$	$\theta(3^N)$
Row 1	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
Row 2	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
Row 3	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>



Shawn sees the example in lecture of a function with amortized runtime and decides to write his own. He claims his implementation of `addFirst()` has a $\theta(1)$ amortized runtime. Is Shawn correct?

* 2 points

```
public void addFirst(int x) {  
    if (size == items.length) {  
        resize(size + 1000);  
    }  
    items[size] = x;  
    size += 1;  
}
```

- ☐ Shawn is correct. Because the calls to `resize()` are very rare (1 out of every 1000 `addFirst()` calls), we can say that his `addFirst()` runs in $\theta(1)$ time on average.
- ☐ Shawn is not correct. For every N calls to `addFirst`, there will be $N/1000$ — or $\theta(N)$ — calls to `resize()`. Since each call to `resize()` takes linear time, the total cost of resizing after N `addFirst`'s will be $\theta(N^2)$ (or $\theta(N)$ per call on average)

What is the runtime of the following method? *

1 point

```
public void eatThinMints(int n) {  
    if (n == 0) return;  
  
    System.out.println("om nom nom");  
  
    for (int i = 0; i < 2; i++) {  
        eatThinMints(n-1);  
    }  
}
```

- ☐ $\theta(\log N)$
- ☐ $\theta(N)$
- ☐ $\theta(N^2)$
- ☐ $\theta(2^N)$



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